Introduction:

Compilation

Sections 1.2 and 1.3

Compiler

- **Input:** Source program
- **Output:** Target program

- **Method:**
- A compiler reads a *source program* (in source language) and translates them into *target program* (in target language).
- Source and target programs are *functionally equivalent*.

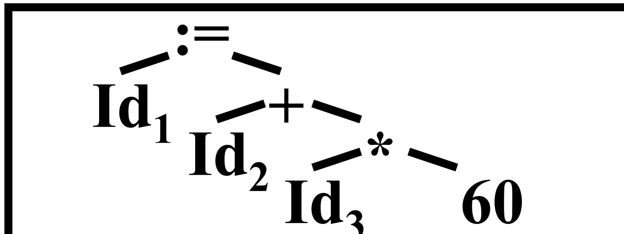
Structure of Compiler: Phases

```
Position := Initial + Rate * 60
```

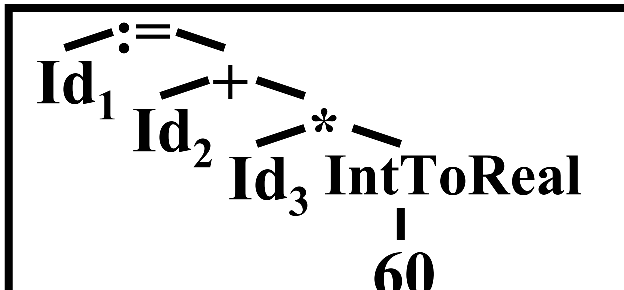
Lexical analyzer

$Id_1 := Id_2 + Id_3 * 60$

Syntax analyzer



Semantic analyzer



Intermediate code generator

```
T1   := IntToReal(60)
T2   := Id3 * T1
T3   := Id2 + T2
Id1  := T3
```

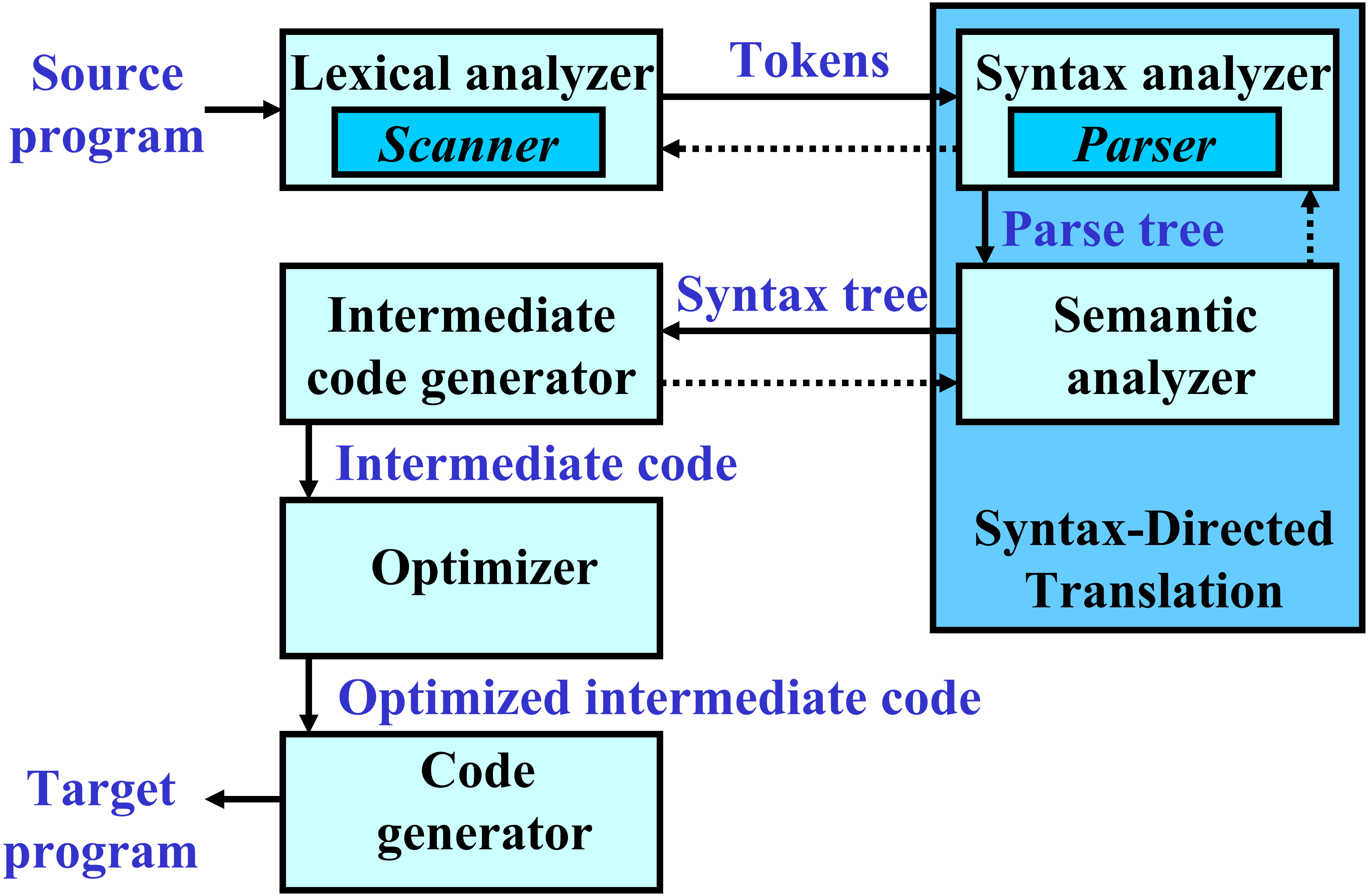
Optimizer

```
T1   := Id3 * 60.0
Id1  := Id2 + T1
```

Code generator

```
fmov R2  , Id3
fmul R2  , #60.0
fmov R3  , Id2
fadd R2  , R3
fmov Id1 , R2
```

Structure of Compiler: Construction



Languages and Compilers

Theoretical view.

$\Sigma = \{a, b\}$, $L = \{a^n b^n : n \geq 0\}$

Question: $aabb \in L$?

Practical view.

Σ = {***begin***, ***end***, ***id***, **:=**, *, **;** , ...},

L_{Pascal} = Programming Language Pascal

Question: ***begin id := id * id; end;*** $\in L_{Pascal}$?

YES: Program is **OK** ⇒ Create a target program

NO: Program is **not OK** ⇒ Handle the errors

Lexical analyzer (Scanner)

- **Input:** Source program
- **Output:** String of tokens

- **Method:**
- Source program is broken into *lexemes* = logically cohesive lexical entities – (identifiers, numbers, key-words, operators,...)
- Lexemes are represented by uniform *tokens*
- Some tokens have *attributes*

Lexical analyzer: Example

Source program:

```
Position := Initial + Rate * 60
```

Lexemes:

Position	:=	Initial	+	Rate	*	60

Tokens:

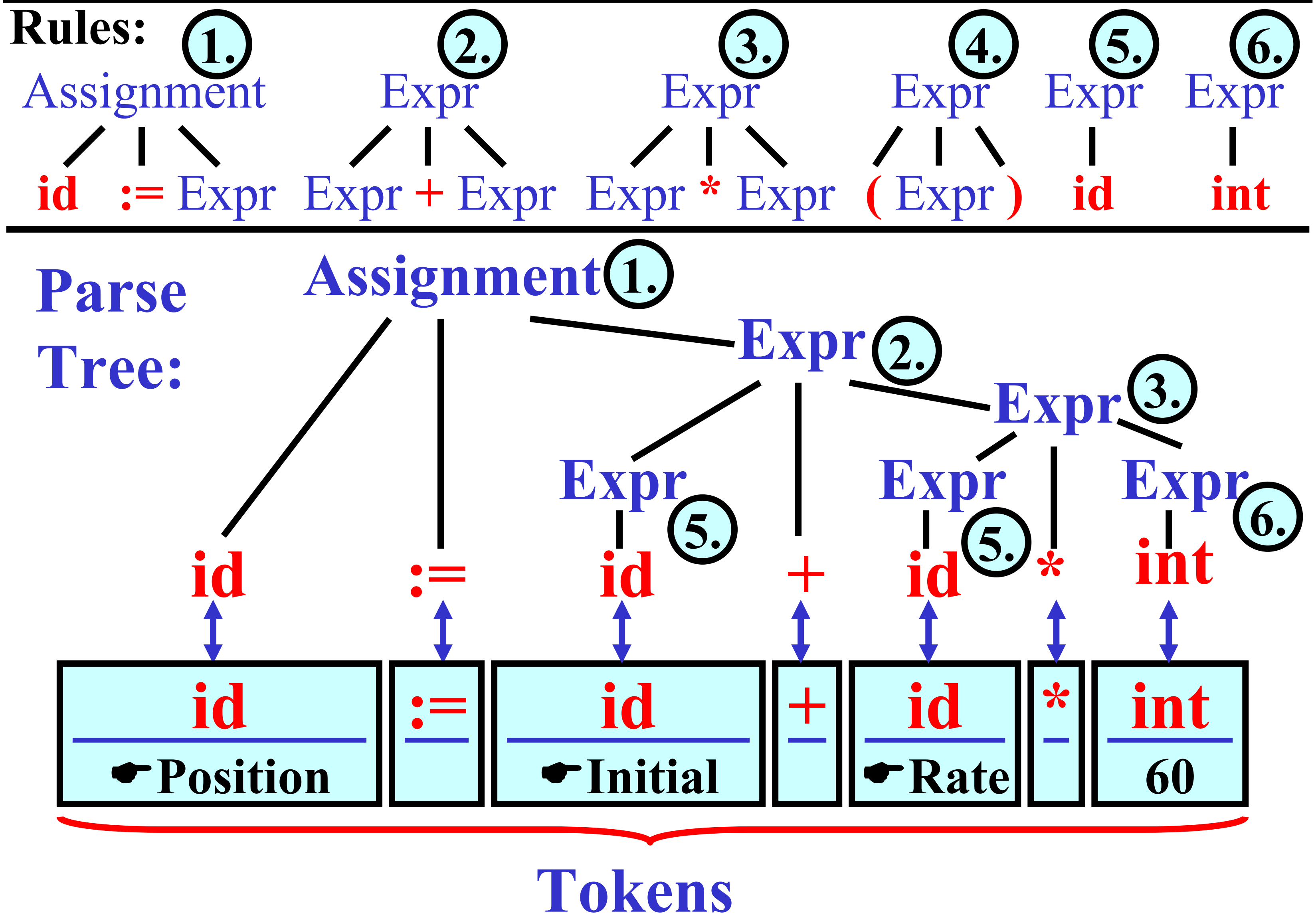
id	:=	id	+	id	*	int
☛Position		☛Initial		☛Rate		60

Syntax analyzer (Parser)

- **Input:** String of tokens
- **Output:** Parse tree

- **Method:**
- Parser verifies that the string of tokens represents a syntactically well-formed program
- If it finds a *parse tree* for the string, it is correct; otherwise, it is not
- Construction of tree is based on grammatical rules
- Two approach: top-down and bottom-up

Syntax analyzer: Example



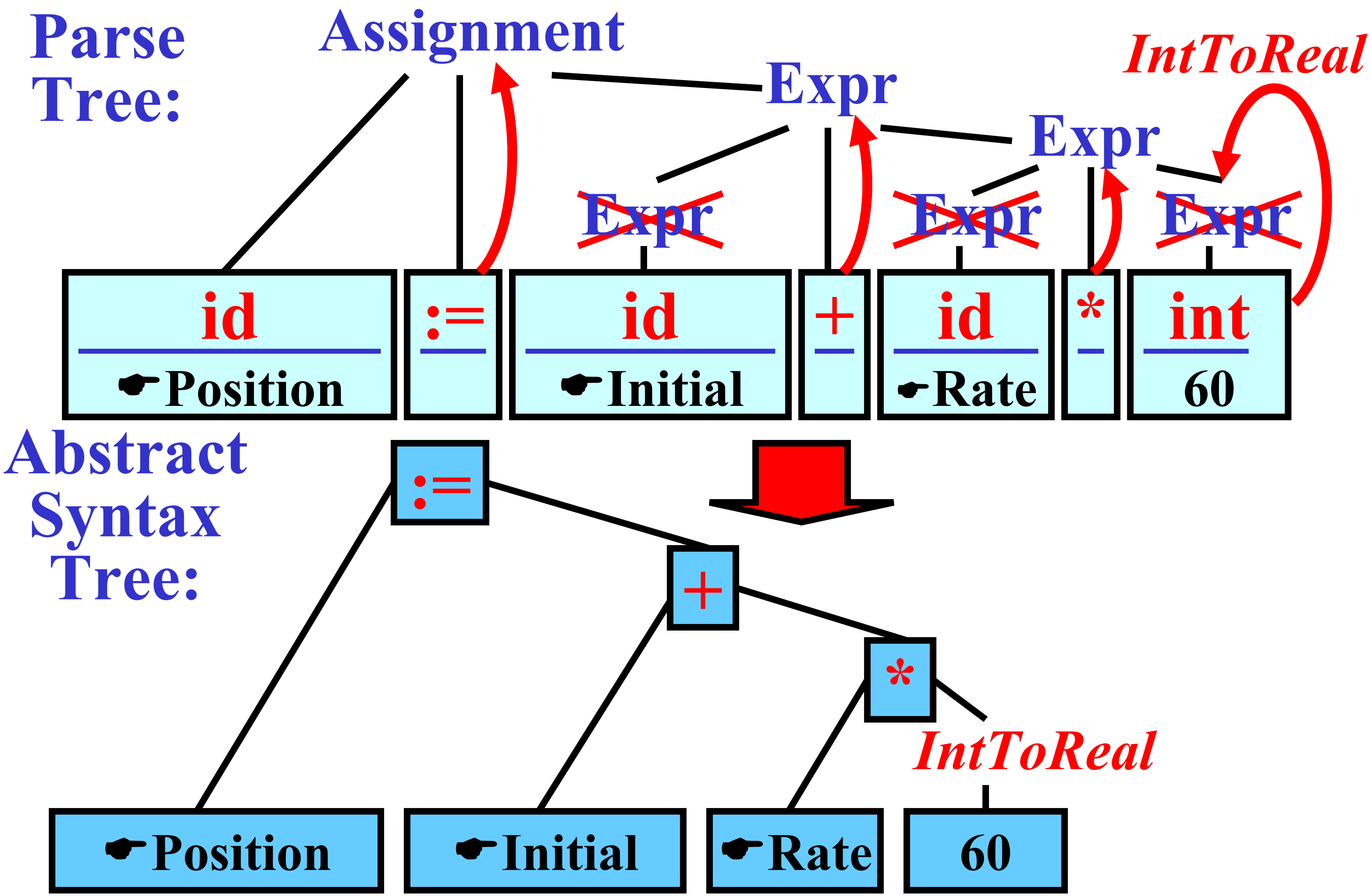
Semantic analyzer

- **Input:** Parse tree
- **Output:** Abstract syntax tree

- **Method:**
- Semantic analyzer checks semantic aspects:
  - *type checking*, which may imply conversions (for example int-to-real)
  - *checking declaration of variables*
- **Syntax-Directed Translation:**

  Parser controls:
  - Semantic actions
  - Generation of syntax tree

Syntax-Directed Translation: Example



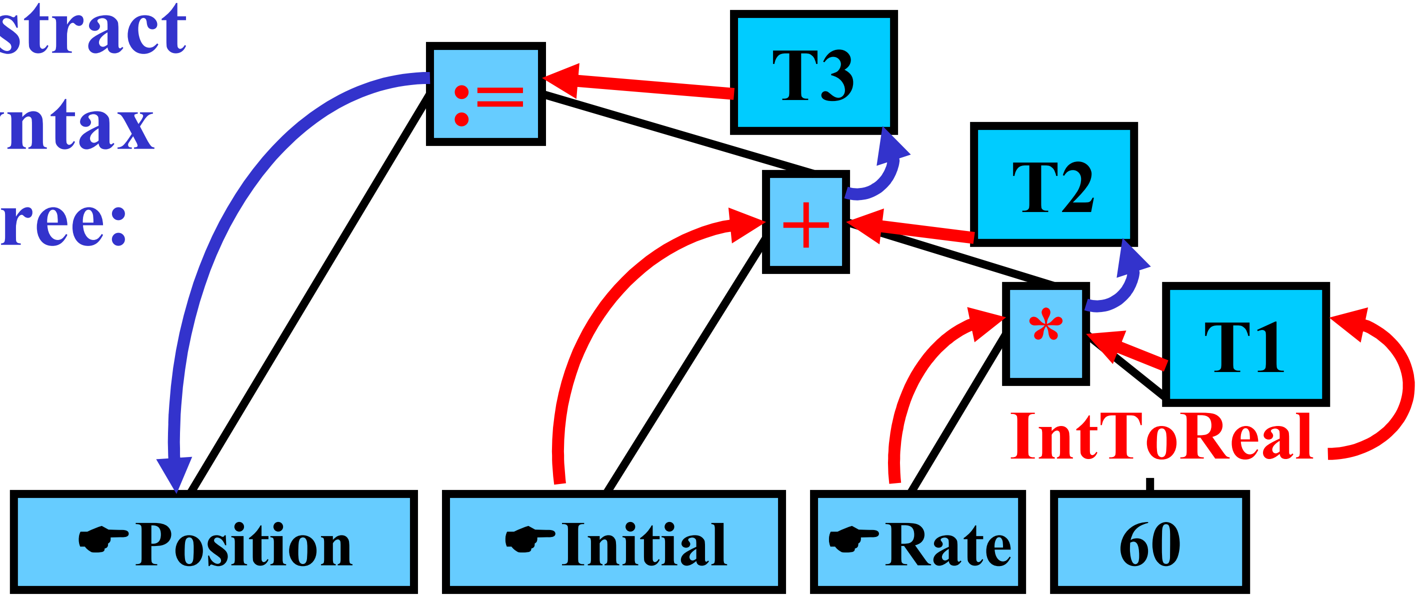
Intermediate code generator

- **Input:** Abstract syntax tree
- **Output:** Intermediate code

- **Method:**
- Intermediate code generator produces the internal version of target program called *intermediate code* for these reasons:
  - uniformity
  - direct translation to target program is difficult and "rough"
  - optimization

Intermediate code generator: Example

Abstract Syntax Tree:



Intermediate code:

```
T1  := IntToReal(60)
T2  := ☛Rate * T1
T3  := ☛Initial + T2
☛Positon := T3
```


Optimizer

- **Input:** Intermediate code
- **Output:** Optimized intermediate code

- **Method:**
- Optimizer makes more efficient version of intermediate code called *optimized intermediate code*:
    - **Constant propagation:** (a := 1; b := 2; c := a + b ⇒ c := 3)
      Note: *Variables a, b have no next use*
    - **Copy propagation:** (b := a; c := b; d := c ⇒ d := a)
      Note: *Variables b, c have no next use*
    - **Dead code elimination:** (**while** false **do** … ⇒ nothing)

⋮

Note: Some compilers have no optimizer

Optimizer: Example

Intermediate code:

```
T1 := IntToReal(60)
T2 := ☛Rate * T1
T3 := ☛Initial + T2
☛Positon:= T3
```

Optimized intermediate code:

```
T2 := ☛Rate * 60.0
☛Positon := ☛Initial+T2
```

Code Generator

- **Input:** Optimized intermediate code
- **Output:** Target program

- **Method:**
- Optimized intermediate code is converted to *target program*
- Target program is written in target language
- In reality, target language is assembly or machine language

Code Generator: Example

Optimized intermediate code:

```
T2 := ☛Rate * 60.0
☛Positon := ☛Initial+T2
```

Target program:

```
fmov R2, ☛Rate
fmul R2, #60.0
fmov R3, ☛Initial
fadd R2, R3
fmov ☛Position,R2
```

R2 ≅ T2