

# Augmented Transition Networks

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- **Introduction**
- **Recursive Transition Network**
- **Augmented Transition Networks**
- **ATN with Global Register**

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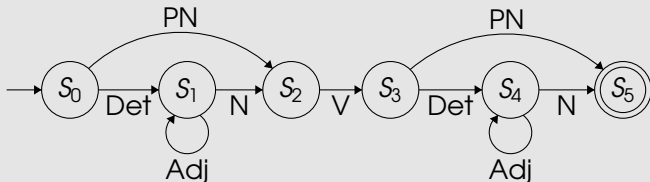
## Augmented Transition Networks (ATN)

- One of the formal description of NL.
- One of the earliest implementations of a system for parsing natural languages.
- Work as finite state machines.

## Definition

- Non-deterministic finite state acceptor (NDA) - quadruple  $(Q, \delta, q_0, F)$  over alphabet  $X$ .

For a simple fragment of English can be constructed in following way:



- where:
  - $X$  ... input vocabulary
  - $C$  ... category vocabulary, made up of grammatical categories of NL (NP – noun phrase, VP – verb phrase, V – verb)
  - function  $CAT : X \rightarrow C$ , such that  $CAT(x)$  is the defined category of  $x$ .

## Example

For example if

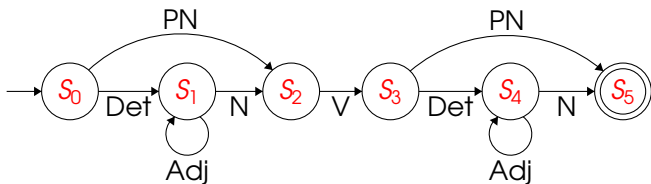
$$X = (\textit{dog}, \textit{John}, \textit{Mary}, \textit{the}, \textit{a}, \textit{loves}, \textit{saw}, \textit{white})$$

and  $C$  is:

$$C = N, \text{Det}, V, \text{Adj}, \text{PN}$$

Using the CAT function we define the following relationships:

- $\text{CAT}(\textit{dog}) = N$
- $\text{CAT}(\textit{John}) = \text{PN}$
- $\text{CAT}(\textit{Mary}) = \text{PN}$
- $\text{CAT}(\textit{the}) = \text{Det}$
- $\text{CAT}(\textit{a}) = \text{Det}$
- $\text{CAT}(\textit{loves}) = V$
- $\text{CAT}(\textit{saw}) = V$
- $\text{CAT}(\textit{white}) = \text{Adj}$



## Example

State  $\{S_i\}$  is created as follows:

- $S_0$  ... initial state
- $S_1$  ... the state in which is legal for the next word to be the noun of the first NP or its adjectives.
- $S_2$  ... the state reached after reading the first NP.
- $S_3$  ... the state reached after reading of a verb.
- $S_4$  ... the state in which is legal for the next word to be adjective or the noun of the next NP.
- $S_5$  ... **completion state** and also  $F = S_5$

## Example

The preceding state diagram corresponds to the following transition function:

- $\delta(S_0, \text{Det}) = \{S_1\}$
- $\delta(S_0, \text{PN}) = \{S_2\}$
- $\delta(S_1, \text{N}) = \{S_2\}$
- $\delta(S_1, \text{Adj}) = \{S_1\}$
- $\delta(S_2, \text{V}) = \{S_3\}$
- for any pair  $(S, x)$  not listed  $\delta(S, x) = \{S_5\}$
- $\delta(S_3, \text{Det}) = \{S_4\}$
- $\delta(S_4, \text{N}) = \{S_5\}$
- $\delta(S_4, \text{Adj}) = \{S_4\}$
- $\delta(S_3, \text{PN}) = \{S_5\}$

The language accepted by the NDA:

$$L = \{\text{PN} \cup \text{Det Adj} * \text{N}\} \cdot \{\text{V}\} \cdot \{\text{PN} \cup \text{Det Adj} * \text{N}\}$$

Example of accepted sentences:

- 1 *John loves Mary.*
- 2 *The white cat saw Mary.*
- 3 *Mary loves a cat.*



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## Problems of NDA formalism

NDA model is unsatisfying:

- The transitions **do not capture the effects of constituency** (there is no way of stating formally that a path from  $S_0$  to  $S_2$  represents a noun phrase)
- FSM **fails to capture** a significant **generalization** in the structure of the language.
- Any kind of **movement dependencies are impossible to capture**.

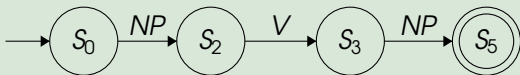
## Problems of NDA model - Solution

- **Allow state transitions** to refer to nonterminals (phrase structures) as well as terminals.
- **Allow a transition between two consecutive states** in RTN to be made by reading **string** (with NDA is possible to read just one word of category  $A$  – string  $w$  belongs to a set  $A$ )
- How do we know that  $w$  belongs to  $A$ ?
  - Because  $w$  has been accepted by *another* RTN – that is why they are called **recursive**.
- **Recursivity** – for recognizing  $A$  we call  $B$  which calls  $C$  etc. until we come to a transition which calls  $A$ -recognizer again.

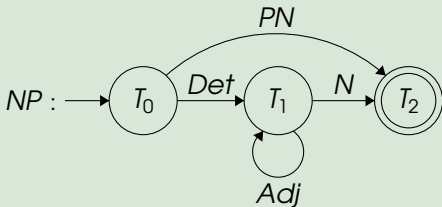
## Example

- The first example after restructuring

The first RTN:



and the second RTN:



## Definition

A **recursive transition network**  $R$  is given by the following data:

- 1  $X$  is a finite set of *terminal* symbols (eg. words of natural language).
- 2  $C$  is a finite set of *categories*.
- 3  $CAT$  is a map  $X \rightarrow C$ .
- 4  $V$  is a set of *nonterminal* symbols.
- 5  $S \in V$  is the *start* symbol.

## Recursive Transition Network

- For each  $v \in V$ , there is a NDA  $M_v = \{Q_v, q_v, \delta_v, F_v\}$ , where  $q_v \in Q_v, F_v \subset Q_v$ , and  $\delta_v : Q_v \times (C \cup V) \rightarrow 2^Q$
- We associate each  $v \in V$  with an **acceptance set**  $T_v$  defined recursively as follows:
  - A string  $w$  is in  $T_v$  just in case it can be partitioned into a sequence  $w_1 w_2 \dots w_n$  such that either  $CAT(w_j) = a_j \in C_j$  or  $w_j \in T_{a_j}$  for  $a_j \in V$  and such that

$$\delta_v^*(q_v, a_1 a_2 \dots a_n) \in F_v.$$

- It means that a string is accepted by  $M_v$  with the transitions from states  $q$  to  $q'$  in  $M_v$  either made with a single  $x$  with  $CAT(x) = c$  such that  $q' \in \delta_v(q, c)$  or by a substring  $w'$  which is accepted by some  $M_u$  for which  $q' \in \delta_v(q, u)$ .
- The language  $L(R)$  accepted by  $R$  is just  $T_S$ , the acceptance set for the start symbol  $S$ .



## Theorem

*For every recursive transition network  $R$  there is a pushdown automaton (PDA)  $P(R)$  which accepts the same language.*



## Problems of RTN model

- There is no way to relate *wh*-sentence with its indicative form.
- Some local independencies are impossible to express (subject-verb agreement).

## Solution

- RTN can be extended by adding **registers** to the model.
- Register may hold arbitrary information about input vocabulary.
- Resulting systems are called **augmented transition network (ATN)**.



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Consider two cases involving agreement

- the subject and verb of a sentence
- the determiner and main noun of an NP.

With the RTN (without register) - there is no way to ensure that sentences like the following are prevented from being accepted as a part of the language.

- 1 ☹️ The *boys is* mischievous.
- 2 ☹️ A *girls* leave home.

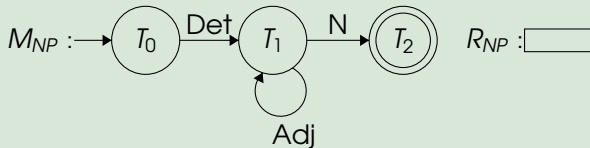
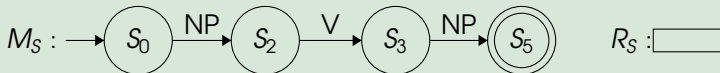
We need to check the number of the subject NP in some way and make the verb transition conditional on the verb agreeing with this NP feature. The following two conditions hold:

Condition I.: If the NP is **singular** then V must be **singular**.

Condition II.: If the NP is **plural** then V must be **plural**.

## Example

Example of what a register can contain and how this information is structured.





## Example

Each register can hold one of three values:

- singular,
- plural,
- neutral.

Singular does not agree with plural, but neutral agrees with every value.

Each NDA is augmented as follows:

Transitions for  $M_S$ , where  $w$  is phrase (or network) variable and  $x$  is a word (or arc) variable:

**Transition 1**  $\delta_S(S_0, w) = S_2$  if  $w \in T_{NP}$ . When transition is made  $R_S := n$ , where  $n$  is value set in  $R_{NP}$  by the acceptance of  $w$  by  $M_{NP}$

**Transition 2**  $\delta_S(S_2, x) = S_3$  if  $CAT(X) = v$  and the number of  $x$  agrees with that stored in  $R_S$ .

**Transition 3**  $\delta_S(S_3, w) = S_5$  if  $w \in T_{NP}$ .



## Example

Transitions for  $M_{NP}$ :

**Transition 4.**  $\delta_{NP}(T_0, x) = T_1$  if  $CAT(X) = Det$ . In making the transition,  $R_{NP} :=$  the number of  $x$ .

**Transition 5.**  $\delta_{NP}(T_1, x) = T_1$  if  $CAT(X) = Adj$  and the number of  $x$  agrees with that stored in  $R_{NP}$ .

**Transition 6.**  $\delta_{NP}(T_2, x) = T_1$  if  $CAT(X) = N$  and the number of  $x$  agrees with that stored in  $R_{NP}$ .

On exit from  $T_2$ , a string  $w$  in  $NP$  carries the number stored in  $R_{NP}$



## Long-distance dependencies

How long-distance dependencies can be handled in an ATN formalism?

Note that neither FSAs nor RTNs were able to parse these sentences (including wh-movement).

- 1 *What does John love?*
- 2 *Who loves Mary?*

With previous 6 transitions we are now **unable** to handle such sentences.

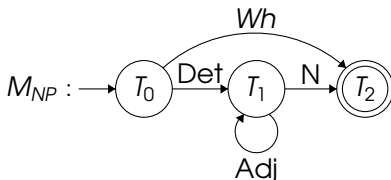
One of the possible solutions is to use **global register**.

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We need some additional rules.

## Global register

- A register which is associated with the whole sentence.
- Any transition can modify and access this register.
- eg. to parse previous sentences  $\rightarrow$  modify ATN in the following way: we add global register  $R_{wh}$  which can be empty, hold *wh* or hold *do*.

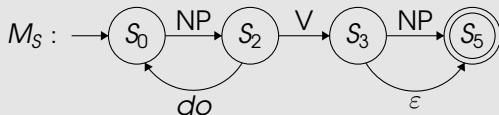


We add new transitions to  $M_{NP}$ :

**Transition 7.**  $\delta_{NP}(T_0, x) = T_2$  if  $CAT(x) = Wh$  in making the transition,  $R_{wh} := x$ , and  $R_{NP} := neutral$ .



And to  $M_S$  we add these two new transitions:



**Transition 8.**  $\delta_M(S_2, x) = S_0$  if  $CAT(X) = do$  and  $R_{wh} := wh$ .  
In making the transition,  $R_{wh} := do$ .

**Transition 9.**  $\delta_M(S_3, \varepsilon) = S_5$  if  $R_{wh} := do$ .

And to Transition 3 we add the condition  $R_{wh} \neq do$

**Transition 3**  $\delta_S(S_3, w) = S_5$  if  $w \in T_{NP}$ ,  $R_{wh} \neq do$ .




This structure accepts both previous sentences 1 and 2 and also sentence 4, but rejects 3 (which is not correct.)

- 1 What does John love?
- 2 Who loves Mary?
- 3 ☹️ Who likes?
- 4 Who likes who?

The ability to augment a transition with arbitrary tests and computation gives an ATN the computational power of **Turing machines**.

## Theorem

*For each Turing machine  $T$ , there is an ATN which accepts precisely the string accepted by  $T$ .*

-  James Allen:  
*Natural Language Understanding*,  
The Benjamin/Cummings Publishing Company. Inc., 2005
-  Robert N. Moll, Michael A. Arbib, A. J. Kfoury:  
*An Introduction to Formal Language Theory*,  
Springer-Verlag, 1988