Regulated Pushdown Automata

Alexander Meduna

Faculty of Information Technology
Brno University of Technology
Brno, Czech Republic, Europe

Fundamental References

- **Meduna Alexander, Kolář Dušan:**
Regulated Pushdown Automata, ***Acta Cybernetica*, Vol. 2000, No. 4, p. 653-664**

- **Meduna Alexander, Kolář Dušan:**
One-Turn Regulated Pushdown Automata and Their Reduction, *Fundamenta Informatica*, Vol. 2002, No. 16, p. 399-405

Inspiration: Regulated Grammars

- **Grammar *G*:**

1. $S \rightarrow AC$

2. $A \rightarrow aAb$

3. $A \rightarrow ab$

4. $C \rightarrow Cc$

5. $C \rightarrow c$

- **$\Xi = \{1\}\{24\}^*\{35\}$**

Regulated Grammars 1/2

- **Grammar *G*:**

1. $S \rightarrow AC$
2. $A \rightarrow aAb$
3. $A \rightarrow ab$
4. $C \rightarrow Cc$
5. $C \rightarrow c$

$\Xi = \{1\}\{24\}^*\{35\}$

- **Without Ξ, *G* generates *aabbccc*:**

$$\begin{aligned} S &\Rightarrow AC & [1] \\ &\Rightarrow aAbC & [2] \\ &\Rightarrow aAbCc & [4] \\ &\Rightarrow aabbCc & [3] \\ &\Rightarrow aabbCcc & [4] \\ &\Rightarrow aabbccc & [5] \end{aligned}$$

$$\boxed{L(G) = \{a^n b^n c^m : n, m \geq 1\}}$$

Regulated Grammars 2/2

- with Ξ, G does not generate ***aabbccc***, because

$$124345 \notin \Xi = \{1\}\{24\}^*\{35\}$$

- with Ξ, G generates ***aabbcc***:

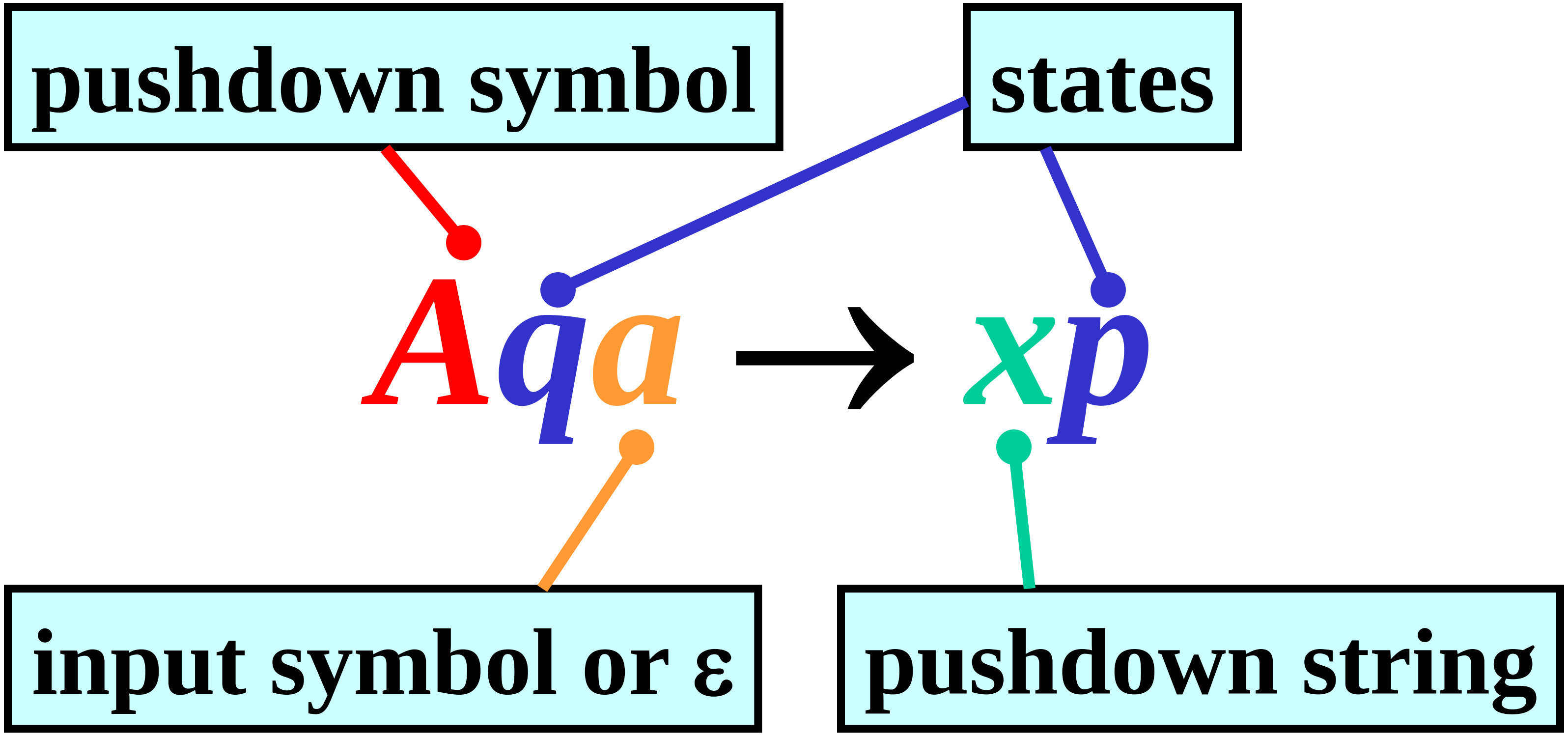
$$\begin{array}{rll}
S & \Rightarrow AC & [1] \\
 & \Rightarrow aAbC & [2] \\
 & \Rightarrow aAbCc & [4] \\
 & \Rightarrow aabbCc & [3] \\
 & \Rightarrow aabbcc & [5]
\end{array}$$

and $12435 \in \Xi$

$$\boxed{L(G, \Xi) = \{a^n b^n c^n : n \geq 1\}}$$

PDA: Notation

- **A PDA is based on a finite set of rules of the form:**



pushdown symbol

states

$$Aqa \rightarrow xp$$

input symbol or ε

pushdown string

New Concept: Regulated PDAs

- **PDA *M*:**

1. $Ssa \rightarrow Sas$
2. $asa \rightarrow aas$
3. $asb \rightarrow q$
4. $aqb \rightarrow q$
5. $Sqc \rightarrow Sq$
6. $Sqc \rightarrow f$

- $\Xi = \{12^m34^n5^n6: m, n \geq 0\}$

Regulated PDAs 1/2

- **PDA *M*:**

  1. $Ssa \rightarrow Sas$
  2. $asa \rightarrow aas$
  3. $asb \rightarrow q$
  4. $aqb \rightarrow q$
  5. $Sqc \rightarrow Sq$
  6. $Sqc \rightarrow f$

$\Xi = \{12^m34^n5^n6: m, n \geq 0\}$

- **Without Ξ, *M* accepts *aabbccc*:**

$Ssaabbccc$	
$\Rightarrow Sasabbccc$	[1]
$\Rightarrow Saasbbccc$	[2]
$\Rightarrow Saqbccc$	[3]
$\Rightarrow Sqccc$	[4]
$\Rightarrow Sqcc$	[5]
$\Rightarrow Sqc$	[5]
$\Rightarrow f$	[6]

$$L(M) = \{a^nb^nc^m: n, m \geq 1\}$$

Regulated PDAs 2/2

- with Ξ, *M* does not accept ***aabbccc*** because

$$\mathbf{1234556} \notin \Xi = \{\mathbf{12}^{m}\mathbf{34}^{n}\mathbf{5}^{n}\mathbf{6}: m, n \geq 0\}$$

- with Ξ, *M* accepts ***aabbcc***:

$$
\begin{aligned}
\mathit{Ssaabbcc} &\Rightarrow \mathit{Sasabbcc} && [1]\\
&\Rightarrow \mathit{Saasbbcc} && [2]\\
&\Rightarrow \mathit{Saqbcc} && [3]\\
&\Rightarrow \mathit{Sqcc} && [4]\\
&\Rightarrow \mathit{Sqc} && [5]\\
&\Rightarrow f && [6]
\end{aligned}
$$

and **123456** $\in \Xi$

$$\boxed{L(M, \Xi) = \{a^{n}b^{n}c^{n}: n \geq 1\}}$$

Gist: Regulated PDAs

- Consider a pushdown automaton, *M*, and control language, Ξ.
- *M* accepts a string, *x*, if and only if Ξ contains a control string according to which *M* makes a sequence of moves so it reaches a final configuration after reading *x*.

Definition: Regulated PDA 1/4

A pushdown automaton is a 7-tuple

$$M = (Q, \Sigma, \Omega, R, s, S, F)\text{, where}$$

- Q is a *finite set of states*,
- Σ is an *input alphabet*,
- Ω is a *pushdown alphabet*,
- R is a *finite set of rules* of the form:

$$Apa \rightarrow wq\text{, where}$$

$$A \in \Omega, p,q \in Q, a \in \Sigma \cup \{\varepsilon\}, w \in \Omega^*$$

- $s \in Q$ is the *start state*
- $S \in \Omega$ is the *start symbol*
- $F \subseteq Q$ is a set of *final states*

Definition: Regulated PDA 2/4

- Let Ψ be an alphabet of *rule labels*. Let every rule $Apa \rightarrow wq$ be labeled with a unique $\rho \in \Psi$ as

$$\rho.\ Apa \rightarrow wq.$$

- A configuration of M, χ, is any string from $\Omega^*Q\Sigma^*$

- For every $x \in \Omega^*$, $y \in \Sigma^*$, and $\rho.\ Apa \rightarrow wq \in R$, M makes a move from configuration $xApay$ to configuration $xwqy$ according to ρ, written as

$$xApay \Rightarrow xwqy\ [\rho]$$

Definition: Regulated PDA 3/4

- Let χ be any configuration of *M*. *M* makes *zero moves* from χ to χ *according to* ε, written as

$$\chi \Rightarrow^0 \chi\ [\varepsilon]$$

- Let there exist a sequence of configurations $\chi_0, \chi_1, \ldots, \chi_n$ for some $n \geq 1$ such that $\chi_{i-1} \Rightarrow \chi_i\ [\rho_i]$, where $\rho_i \in \Psi$, for $i = 1,\ldots,n$, then *M* makes *n moves from* χ_0 *to* χ_n *according to* $[\rho_1 \ldots \rho_n]$, written as

$$\chi_0 \Rightarrow^n \chi_n\ [\rho_1 \ldots \rho_n]$$

Definition: Regulated PDA 3/4

- If for some $n \geq 0$, $\chi_0 \Rightarrow^n \chi_n\ [\rho_1 \ldots \rho_n]$, we write

$$\chi_0 \Rightarrow^* \chi_n\ [\rho_1 \ldots \rho_n]$$

- Let Ξ be a *control language* over Ψ, that is, $\Xi \subseteq \Psi^*$. With Ξ, M accepts its language, $L(M, \Xi)$, as

$$L(M, \Xi) = \{w\colon w \in \Sigma^*, Ssw \Rightarrow^* f\ [\sigma], \sigma \in \Xi\}$$

Language Families

- ***LIN*** - the family of linear languages
- ***CF*** - the family of context-free languages
- ***RE*** - the family of recursively enumerable languages

- ***RPD(REG)*** - the family of languages accepted by PDAs regulated by regular languages
- ***RPD(LIN)*** - the family of languages accepted by PDAs regulated by linear languages

Theorem 1 and its Proof 1/2

$RPD(REG) = CF$

Proof:

I. $CF \subseteq RPD(REG)$ is clear.

II. $RPD(REG) \subseteq CF$**:**

- Let $L = L(M, \Xi)$,

PDA **Regular language**

- Let $\Xi = L(G)$, G - regular grammar based on rules: $A \rightarrow aB$, $A \rightarrow a$

Theorem 1 and its Proof 2/2

Transform *M* regulated by Ξ to a *PDA N* as follows:

1) for every $a.Cqb \rightarrow xp$ from *M* and
every $A \rightarrow aB$ from *G*,
add $C\langle qA\rangle b \rightarrow x\langle pB\rangle$ to *N*

2) for every $a.Cqb \rightarrow xp$ from *M* and
every $A \rightarrow a$ from *G*,
add $C\langle qA\rangle b \rightarrow x\langle pf\rangle$ to *N*

New symbol

3) The set of final states in *N*:

$\{\langle pf\rangle : p$ is a final state in $M\}$

Theorem 2

$$\boldsymbol{RPD(LIN) = RE}$$

Proof:

- See [**Meduna Alexander, Kolář Dušan: Regulated Pushdown Automata**, ***Acta Cybernetica*****,Vol. 2000, No. 4, p. 653-664**]

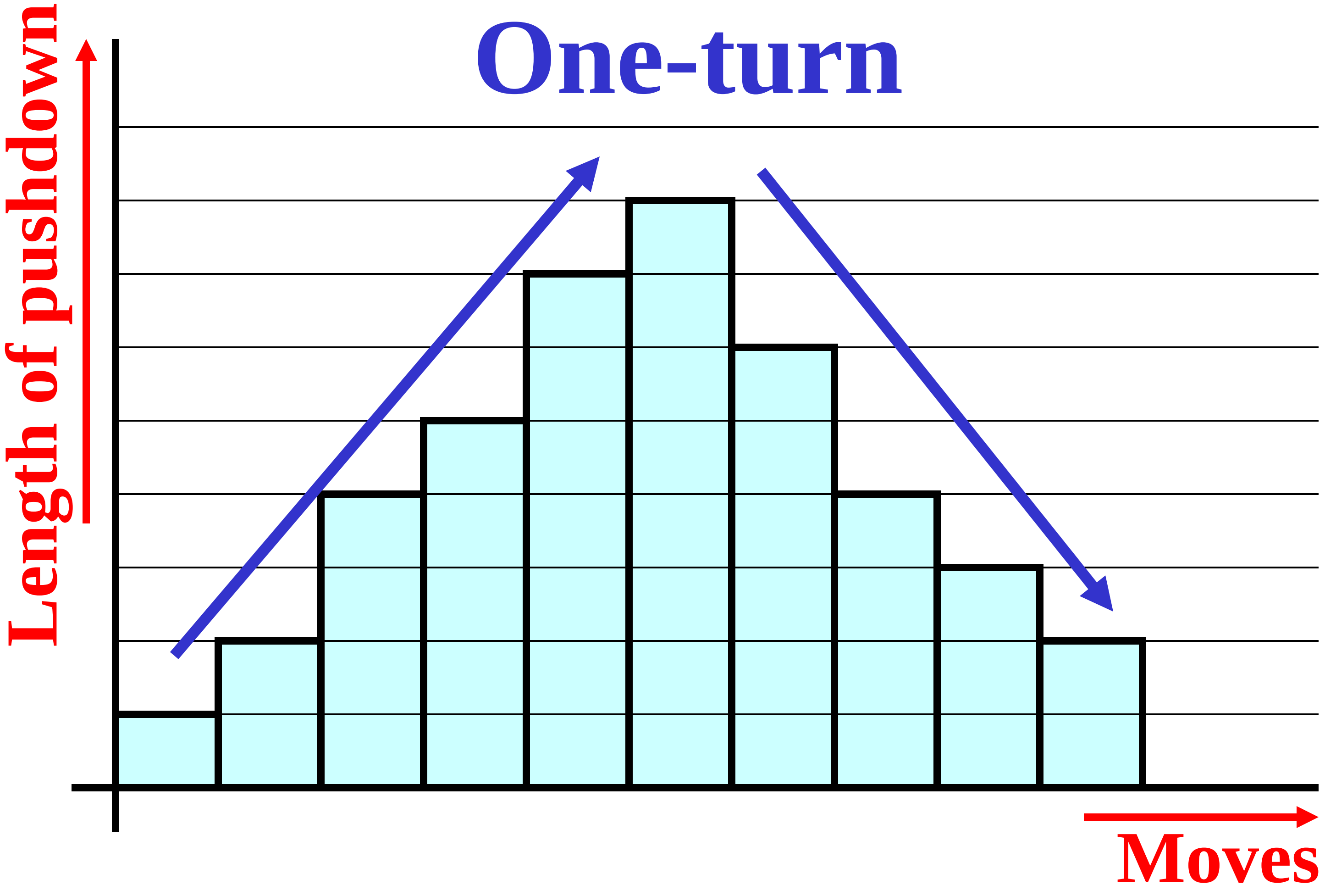
Simplification of RPDAs 1/2

I. consider two consecutive moves made by a pushdown automaton, *M*.

If during the first move *M* does not shorten its pushdown and during the second move it does, then *M* makes ***a turn*** during the second move.

- A pushdown automaton is *one-turn* if it makes no more than one turn during any computation starting from an initial configuration.

One-Turn PDA: Illustration



Simplification of RPDAs 2/2

II. During a move, an *atomic* regulated PDA changes a state and, in addition, performs exactly one of the following actions:

1. pushes a symbol onto the pushdown
2. pops a symbol from the pushdown
3. reads an input symbol

Theorem 3

- **Every $L \in RE$ is accepted by an atomic one-turn PDA regulated by Ξ, where $\Xi \in LIN$.**

Proof:

- See [**Meduna Alexander, Kolář Dušan: One-Turn Regulated Pushdown Automata and Their Reduction, *Fundamenta Informatica*,Vol. 2002, No. 16, p. 399-405**]

End